

An Introduction to Graph Folding

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Abstract

Graph folding is a unary operation introduced by Gervacio [6] in 1992 which he later modified in 1999.

This operation induces in a very natural way a partitioning of the vertex set of a graph into pairwise linked and independent sets. In this paper we give results involving the folding of regular graphs. Here, too, we show that the famous Petersen graph folds only into the complete graphs K_3 , K_4 , and K_5 . In addition, among others, we will show that $\max\{t \mid K_t \in F(G)\} \leq \left\lfloor \frac{1+\sqrt{1+4rn}}{2} \right\rfloor$, where G is a connected r -regular graph of order n .

Furthermore, this paper enumerates some relationships of folding graphs to several graph invariants such as chromatic number, span, independence number, domination number, maximum degree, and dimension.

Keywords: graph, folding, chromatic number, span, dimension

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1 Introduction

We shall only consider connected simple graphs here, i.e., finite graphs without multiple edges and loops. *Folding a graph* simply means identifying two non-adjacent vertices of a graph and this can be repeated iteratively any number of times until there does not exist such pair of non-adjacent vertices.

A master's thesis [12] and a doctoral dissertation [9] were completed and successfully defended in 1995 and 2002, respectively. Furthermore, in August 2001 Gervacio presented a paper on this topic at a Seminar-Lecture series at University of Santo Thomas, Manila, sponsored by the Mathematical Society of the Philippines, National Capital Region Chapter.

In 1992, Gervacio [6] studied folding of graphs using a different construction. In that study, a folding of a connected graph G is defined as the simple graph obtained from G by identifying as one vertex two non-adjacent vertices and then reducing multiple edges formed to single edges. In contrast, this present study defines folding similar to the definition mentioned above with the added condition that two non-adjacent vertices to be identified must have a common neighbor.

Some of the results obtained from Gervacio's work [6] included the relationship of folding graphs with concepts of a perimeter and span of graphs. Formulas for the span and perimeter of a connected graph were also derived.

In her master's thesis, Vega [12] made further studies of folding graphs using the definition used in [6]. Results obtained include, among others, the determination of the *maximum-folding* and *minimum-folding* of a graph G , which are denoted by $mx f(G)$ and $mn f(G)$, respectively. Here $mx f(G)$ is defined as the k -folding of a graph G , where k is the largest non-negative

integer such that the resulting graph is a complete graph of minimum order while $mnf(G)$ is defined as the k -folding of G , where k is the smallest non-negative integer for which the resulting complete graph is of maximum order.

As a matter of comparison, it should be stressed here that the folding used in this study may also be referred to as restricted folding while that in [6] and [12] is a non-restricted or free folding. Notice that in a restricted folding we identify two distinct vertices of a graph whose distance is 2 while in free folding we identify two distinct vertices whose distance is greater than or equal to 2.

Buckley and Harary [1] stated that Moon and Mosher proved that almost all graphs have diameter 2. Hence, under this construction, we may then be able to fold almost all graphs.

Observe further that a restricted folding is a free folding but not conversely. For instance, it will be shown that the cycle C_6 of order 6 restrictively folds uniquely into the complete graph K_2 ; while the same graph can be freely folded into K_3 . In fact, it was shown in [12] that $mnf(C_6) = K_3$. However, $mx f(G) = K_2$ also. Thus, for an arbitrary graph G it is not always true that we can obtain $mnf(G)$ using the folding used in this present study (restricted, as we referred it here).

It should be noted that in the cases of folding the wheel, fan, biwheel and bifan, the same minimum foldings were obtained using these two constructions of folding. This is expected because of the structures of these graphs.

Recently, Gervacio [7, 8] made separate investigations related to this study using restricted folding. Some of the results of these works have relevance

on this study. One result obtained is a characterization theorem for folding graphs. Other important results are the partitioning the vertex-sets of the path and cycle into pairwise linked and independent sets. These results enable us to determine the exact orders of the maximum complete graphs into which the wheel and the fan can be folded.

2 The Idea of Graph Folding, Preliminary Concepts and Results

Unless stated otherwise, graphs will be denoted by capital letters such as G, H , etc. The set of vertices of G and the set of edges of G will be denoted by $V(G)$ and $E(G)$, respectively. Basic graph-theoretic concepts are assumed to be understood by the reader; otherwise, the readers may refer to the books by Chartrand [2], Foulds[3], and Harary [10]. New concepts will be defined formally.

Definition 2.1 Let G be a connected graph and let x and y be non-adjacent vertices with at least one common neighbor. We define the xy -folding of G (or simply the folding of G when no confusion can result), denoted by $f(G; x, y)$ (or, $f(G; y, x)$), to be the graph obtained from G by identifying the vertices x and y and reducing multiple edges to single edges.

Notation: Let x_1, x_2, \dots, x_n be vertices of a graph G . The notation $\{x_1, x_2, \dots\}$ means that x_1, x_2, \dots, x_n are identified.

Example 2.2 The graph G in Figure 1 has two vertices x and y with common neighbors u and v . By identifying x and y and reducing the resulting

double edges into simple edges, we get the graph $f(G, x, y)$. Note that u and z cannot be identified since they do not have a common neighbor.

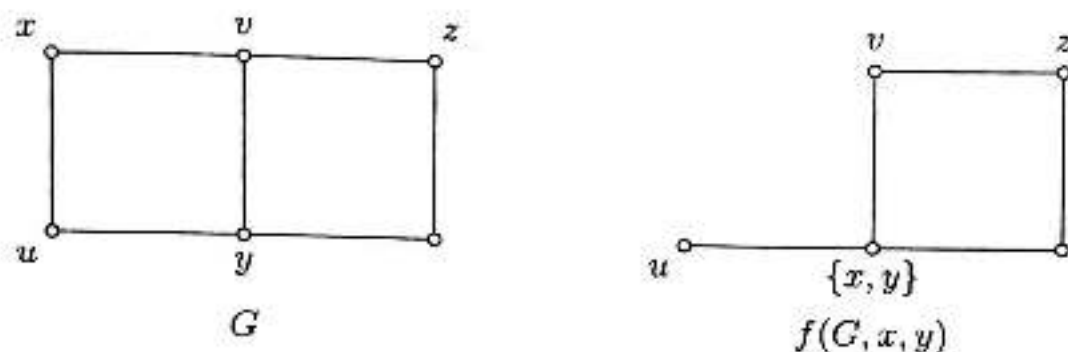


Figure 1: Folding of a Graph G into $f(G, x, y)$.

Observe, further, that if a connected graph is not complete, then it always has two non-adjacent vertices that have a common neighbor. Thus, any connected non-complete graph can undergo a sequence of folding until we obtain a complete graph. Obviously, no loops can occur in any series of folding by definition.

Notation: We shall denote by $F(G)$ the set of all non-isomorphic complete graphs obtainable from a connected graph G by a sequence of folding. That is,

$$F(G) = \{K_p \mid G \text{ can be folded into } K_p\}.$$

If $F(G)$ has only one element, say K_s , then we shall simply write $F(G) = K_s$ instead of $F(G) = \{K_s\}$. For instance, it is obvious that $F(K_n) = K_n$, $\forall n$.

Let us fold further the graph G in Figure 1 until we get a complete graph. (Please refer to Figure 2)

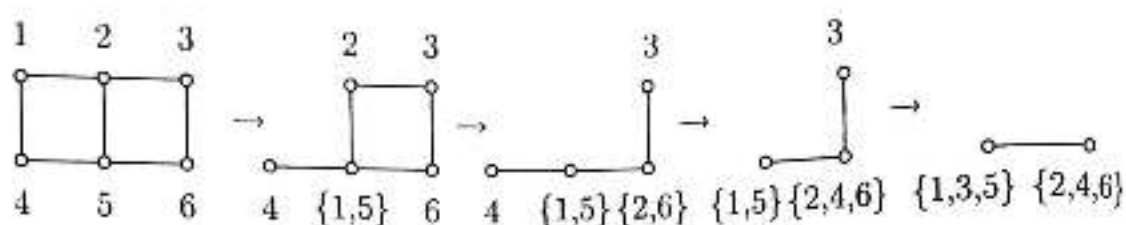


Figure 2: Folding a graph into a complete graph.

The graph was folded into the complete graph K_2 . It was shown in [9] that a connected graph G folds uniquely into K_2 if and only if G is bipartite. In view of this result, we have $F(G) = K_2$ since the graph G in Figure 1 is bipartite.

Let us observe that each of the two vertices of K_2 was obtained by identifying some vertices of the original graph. One of the vertices of K_2 was obtained by identifying the vertices 1, 3 and 5 while the other vertex was obtained by the vertices 2, 4, and 6. Observe that the sets $\{1, 3, 5\}$ and $\{2, 4, 6\}$ are both independent and pairwise linked sets which partitions $V(G)$. This observation is true for all graphs according to [9].

Indeed there are graphs G for which $F(G)$ contains more than one complete graph as can be seen in the following example.

Example 2.3 We show that $F(W_6) = \{K_3, K_4\}$.

Solution: By definition, $W_6 = C_6 + K_1$. In general, W_n is commonly called the *wheel* of order $n + 1$ while C_n is called the *cycle* of order n . The graph of W_6 is shown in Figure 3.

We now fold this graph in two ways:

Case 1. If we choose x and y to be vertices whose distance in C_6 is 3, and then continue folding, we eventually get K_4 .

Case 2. If in the first folding we choose two alternate vertices x and y in C_6 and then continue folding, we eventually get K_3 .

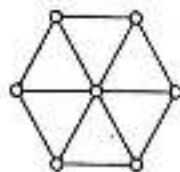


Figure 3: The Wheel W_6

Example 2.4 We invite the interested reader to verify that $F(F_9) = \{K_3, K_4, K_5\}$, where F_9 is the fan of order 9.

We end this section by stating the following results which will prove relevant later on.

Theorem 2.5 [9] *If $K_p \in F(G)$, then $V(G)$ can be partitioned into p independent sets S_1, S_2, \dots, S_p which are pairwise linked.*

Theorem 2.6 [2] *Let G be a graph of order n and $v \in V(G)$. If $\deg_G(v) = r$, then $\deg_{\tilde{G}}(v) = n - r - 1$, where \tilde{G} is the complement of G . Therefore, \tilde{G} is regular if and only if G is regular.*

3 Folding Regular Graphs

This section deals on the folding of some regular graphs.

A graph G is called *regular* if all the vertices of G are of equal degree. If $\forall x \in G, \deg(x) = r$, we say that G is an *r -regular graph*.

First, let us find all the distinct foldings of a very popular graph – the Petersen Graph which we shall denote by G^* . This graph is shown in Figure 3 together with a labelling of its vertices. Note that G^* is 3-regular.

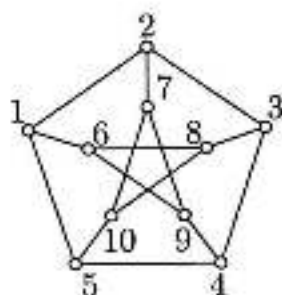


Figure 4: The Petersen graph G^* .

Theorem 3.1 *The Petersen graph G^* folds only into K_3, K_4 , and K_5 , i.e., $F(G^*) = \{K_3, K_4, K_5\}$.*

Proof: Shown in Figure 3 is the Petersen graph with a labelling of its vertices. By a result in [9] mentioned earlier, $K_2 \notin F(G^*)$, since G^* is not bipartite.

A 3-partition of $V(G^*)$ that folds G^* into K_3 is the following:

$$S_1 = \{1, 8, 9\}, S_2 = \{3, 5, 7\}, S_3 = \{2, 4, 6, 10\}.$$

(Please refer to Figure 5.)

A 4-partition of $V(G^*)$ that folds G^* into K_4 is the following:

$$S_1 = \{1, 8, 4\}, S_2 = \{3, 5, 9\}, S_3 = \{2, 10\}, S_4 = \{6, 7\}.$$

A 5-partition of $V(G^*)$ that folds G^* into K_5 is the following:

$$S_1 = \{1, 4\}, S_2 = \{8, 9\}, S_3 = \{3, 10\}, S_4 = \{2, 6\}, S_5 = \{5, 7\}.$$

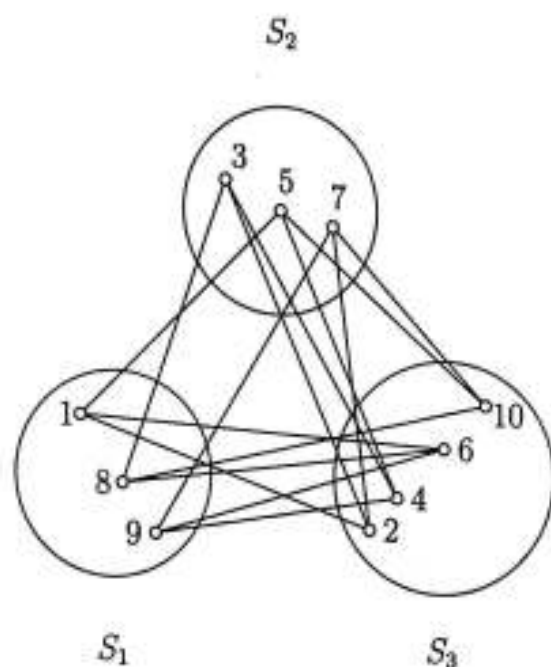


Figure 5: Partitioning $V(G^*)$ that folds G^* into K_3 .

Therefore we succeeded in folding the Petersen graph into K_3 , K_4 and K_5 .

We now claim that G^* cannot be folded into K_6 and, hence into complete graphs of order $n > 6$. To see this, observe that each time we identify two non-adjacent vertices, the number of edges decreases. Note that G^* has 15 edges and K_6 has $\binom{6}{2} = 15$ edges also. Hence, we cannot fold it into K_6 .

We now conclude that $F(G^*) = \{K_3, K_4, K_5\}$. \square

Theorem 3.2 *Let G be a k -regular graph of diameter 2, where $k \geq 3$. If G folds into K_r , $r \geq k + 2$, then G necessarily folds into K_{r-1} . In particular, any 3-regular graph of diameter 2 that folds into K_r , where $r \geq 5$, necessarily folds into K_{r-1} .*

Proof: Let S_1, S_2, \dots, S_r be the independent and pairwise linked sets in-

duced by a folding of the k -regular graph G of diameter 2 into K_r . Let $x \in V(G)$. Without loss of generality, let $x \in S_1$. Then $\deg(x) = k \leq r - 2$. Thus x is adjacent to at most $r - 2$ vertices. But there are $r - 1$ sets S_2, S_3, \dots, S_r and $r - 1 > r - 2 \geq k$. Hence there exists some set S_j , $2 \leq j \leq r$, such that x is not adjacent to any vertex in S_j . We can transfer x to S_j maintaining the independence of S_j . Thus, every vertex in S_1 can be transferred to some set among S_2, S_3, \dots, S_r . This leaves us with $r - 1$ sets S_2, S_3, \dots, S_r that are pairwise linked and independent.

Therefore G can now be folded into K_{r-1} . \square

Theorem 3.3 Let G be a connected r -regular graph of order n . Then

$$\max\{t \mid K_t \in F(G)\} \leq \left\lfloor \frac{1 + \sqrt{1 + 4rn}}{2} \right\rfloor.$$

Proof: The theorem is clearly true when $G = K_n$. So assume G is an r -regular graph, where $r \neq n - 1$. Then, $\forall x \in V(G)$, $\deg(x) = r$.

Let $p = \max\{t \mid K_t \in F(G)\}$. By Theorem 2.5, $V(G)$ can be partitioned into p -independent and pairwise linked sets S_1, S_2, \dots, S_p . Without loss of generality, assume

$$k = |S_1| \leq |S_2| \leq \dots \leq |S_p|.$$

Since $\deg(x) = r$, $\forall x \in S_1$ and S_1 is linked to all the other $p - 1$ sets S_2, S_3, \dots, S_p , then,

$$rk \geq p - 1.$$

Furthermore, since the S_i 's partition $V(G)$, then,

$$n \geq pk.$$

From the last two inequalities above, we obtain

$$n \geq p \binom{p-1}{k}$$

by eliminating k . Thus

$$\begin{aligned} nr &\geq p^2 - p \\ p^2 - p - nr &\leq 0 \\ p &\leq \frac{1 + \sqrt{4rn + 1}}{2} \\ p &\leq \left\lfloor \frac{1 + \sqrt{1 + 4rn}}{2} \right\rfloor. \end{aligned}$$

Therefore

$$\max\{t \mid K_t \in F(G)\} \leq \left\lfloor \frac{1 + \sqrt{1 + 4rn}}{2} \right\rfloor. \quad \square$$

Theorem 3.4 *Let G be a connected r -regular graph of order n and such that \tilde{G} is also connected. Then*

$$\max\{t \mid K_t \in F(\tilde{G})\} \leq \left\lfloor \frac{1 + \sqrt{1 + 4n(n-r-1)}}{2} \right\rfloor.$$

Proof: This follows from Theorems 3.3 and 2.6. \square

4 Other Concepts and Initial Results

Here we introduce some relevant concepts as well as results obtained in [9] and other related works.

Other concepts which are not explicitly defined here may be found in [10] and [2].

Let G be a graph.

The largest of the vertex degrees of G is called the *maximum degree* of G and is denoted by $\Delta(G)$ or simply Δ . A *coloring* of G is an assignment of colors to its vertices (exactly one color for each vertex) so that no two adjacent vertices in G have the same color. A coloring of G which assigns c colors to its vertices is called a c -*coloring*.

The *chromatic number* of G is the minimum number, denoted by $\chi(G)$, for which G has a χ -coloring. Equivalently, $\chi(G)$ is the minimum number n for which G is an n -partite graph. Graph G is c -colorable if $\chi(G)$ is less than or equal to c . In other words, if there exists a c -coloring of a graph G , then G is c -colorable. A *clique* in G is a maximal complete subgraph of G . The maximum order of a clique is called the *clique number* of G , and is denoted by $\omega(G)$.

A set S of vertices in G is *independent* if no two vertices of S are adjacent in G . The maximum cardinality of an independent set of vertices in G is called the *independence number* of G and is denoted by $\alpha(G)$. S is a *dominating set* if every vertex not in S is adjacent to a vertex in S . The *domination number* $\gamma(G)$ is the minimum cardinality of a dominating set in G .

G is called a *unit graph* in the Euclidean space R^d if there is a one-to-one mapping $\phi: V(G) \rightarrow R^d$ such that $|\phi(x) - \phi(y)| = 1$ whenever $[x, y] \in E(G)$. The mapping ϕ is called a *unit representation* of G in R^d . The *dimension* of G , denoted by $\dim(G)$, is the smallest integer d such that G is a unit graph in R^d .

Let G be a unit graph in R^d . The *span* of G in R^d , denoted by $\text{span}_d(G)$,

is the real number s such that for every $\epsilon > 0$, the following conditions are satisfied:

1. There exists a unit representation of G in R^d which is contained in some open ball of diameter $s + \epsilon$;
2. No unit representation of G in R^d is contained in any open ball of diameter s .

Let S and T be two disjoint independent sets of vertices of a graph. We say that S and T are *linked* if there is a vertex in S that is adjacent to a vertex in T .

Remark 4.1 [9] *If $K_p \in F(G)$, i.e., if G can be folded into K_p , then G is p -partite.*

Lemma 4.2 [2] *For every graph G , $\gamma(G) \leq \alpha(G)$.*

Lemma 4.3 [3] *If G is a graph of order n and chromatic number $\chi(G)$. Then*

$$\chi(G)\alpha(G) \geq n.$$

Lemma 4.4 [2] *Let G be a connected graph with maximum degree $\Delta = \Delta(G)$. Then*

1. $\omega(G) \leq \chi(G) \leq 1 + \Delta$;
2. $\chi(G) \leq \Delta$ if and only if G is neither a complete graph nor an odd cycle.

Part (2) is called Brook's Theorem.

Lemma 4.5 [5] *For any graph G ,*

$$\dim(G) \leq 2\chi(G).$$

Theorem 4.6 [11] *For any graph G with chromatic number $\chi(G)$,*

$$\text{span}_{2\chi(G)}(G) \leq \sqrt{\frac{2[\chi(G) - 1]}{\chi(G)}}.$$

5 Relation of Folding with Some Graph Invariants

Let $K_p \in F(G)$. By Lemma 4.3, $\alpha(G)\chi(G) \geq n$. Hence

$$p \leq n \leq \alpha(G)\chi(G).$$

Therefore

$$p \leq \alpha(G)\chi(G).$$

This establishes the following:

Remark 5.1 Let G be a graph of order n with independence number $\alpha(G)$ and chromatic number $\chi(G)$. If $K_p \in F(G)$, then

$$p \leq \alpha(G)\chi(G).$$

Note that equality holds when $G = K_n$; for then $\chi(G) = n$, $\alpha(G) = 1$, and $n = p$.

Theorem 5.2 *Let G be a graph with chromatic number $\chi(G)$. Then*

$$\min\{r \mid K_r \in F(G)\} \geq \chi(G).$$

Proof: Let $p = \min\{r \mid K_r \in F(G)\}$. Then $K_p \in F(G)$. Thus G is a p -partite graph by Remark 4.1. Let S_1, S_2, \dots, S_p be a p -partition of $V(G)$. By Theorem 2.5, these sets are pairwise linked sets. Then G is p -colorable since each set S_i may receive one color different from the others. Note that in getting the chromatic number $\chi(G)$, we actually partitioned $V(G)$ into independent sets $S_1, S_2, \dots, S_{\chi(G)}$ which are not necessarily mutually linked.

It follows that $p \geq \chi(G)$, by definition of $\chi(G)$.

Therefore

$$\min\{r \mid K_r \in F(G)\} \geq \chi(G). \quad \square$$

Corollary 5.3 *Let G be a graph with chromatic number $\chi(G)$. Then*

$$\min\{r \mid K_r \in F(G)\} \geq \left\lceil \frac{2}{2 - [\text{span}_{2\chi(G)}(G)]^2} \right\rceil,$$

where $2 \neq [\text{span}_{2\chi(G)}(G)]^2$.

Proof: Let $p = \min\{r \mid K_r \in F(G)\}$. Using Theorem 4.6, we have the following

$$\text{span}_{2\chi(G)}(G) \leq \sqrt{\frac{2[\chi(G) - 1]}{\chi(G)}}$$

$$\chi(G)[\text{span}_{2\chi(G)}(G)]^2 \leq 2\chi(G) - 2$$

$$2 \leq \chi(G) \{2 - [\text{span}_{2\chi(G)}(G)]^2\}$$

$$\chi(G) \geq \frac{2}{2 - [\text{span}_{2\chi(G)}(G)]^2}$$

By Theorem 5.2,

$$p = \min\{r \mid K_r \in F(G)\} \geq \frac{2}{2 - [\text{span}_{2\chi(G)}(G)]^2}$$

Hence,

$$p = \min\{r \mid K_r \in F(G)\} \geq \left\lceil \frac{2}{2 - [\text{span}_{2\chi(G)}(G)]^2} \right\rceil$$

This completes the proof. \square

Corollary 5.4 *Let G be a graph with clique number $\omega(G)$. Then*

$$\min\{r \mid K_r \in F(G)\} \geq \omega(G).$$

Proof: Let $p = \min\{r \mid K_r \in F(G)\}$. By Theorem 5.2, $p \geq \chi(G)$. By Lemma 4.4, $\chi(G) \geq \omega(G)$. Thus, $p \geq \omega(G)$.

Therefore

$$\min\{r \mid K_r \in F(G)\} \geq \omega(G). \quad \square$$

Corollary 5.4 implies that $K_{\omega(G)}$ is a subgraph of K_p .

Theorem 5.5 *Let G be a graph of order n with domination number $\gamma(G)$ and independence number $\alpha(G)$. Then*

$$\min\{r \mid K_r \in F(G)\} > \left\lceil \frac{\gamma(G)}{\alpha(G)} \right\rceil.$$

Proof: By Lemma 4.2,

$$\gamma(G) \leq \alpha(G).$$

Thus

$$\frac{\gamma(G)}{\alpha(G)} \leq 1.$$

Clearly

$$\frac{\gamma(G)}{\alpha(G)} > 0.$$

Therefore

$$0 < \frac{\gamma(G)}{\alpha(G)} \leq 1.$$

But $\min\{r \mid K_r \in F(G)\} \geq 2$; thus,

$$\min\{r \mid K_r \in F(G)\} > \frac{\gamma(G)}{\alpha(G)}.$$

$$\min\{r \mid K_r \in F(G)\} > \left\lceil \frac{\gamma(G)}{\alpha(G)} \right\rceil. \quad \square$$

Theorem 5.6 *Let G be a graph of order n with independence number $\alpha(G)$. Then*

$$\min\{r \mid K_r \in F(G)\} \geq \left\lceil \frac{n}{\alpha(G)} \right\rceil.$$

Proof: Let $p = \min\{r \mid K_r \in F(G)\}$. By Remark 4.1, G is p -partite. Thus, $V(G)$ can be partitioned into p independent sets S_1, S_2, \dots, S_p . Thus we have the following disjoint union,

$$V(G) = S_1 \cup S_2 \cup \dots \cup S_p.$$

Therefore

$$|V(G)| = n = |S_1| + |S_2| + \dots + |S_p|.$$

Since $|S_i| \leq \alpha(G), \forall i$, then

$$n \leq \alpha(G) + \alpha(G) + \dots + \alpha(G),$$

a sum of p terms.

Hence $n \leq p \alpha(G)$, i.e., $p \geq \frac{n}{\alpha(G)}$. Since p is a positive integer, we obtain

$$p \geq \left\lceil \frac{n}{\alpha(G)} \right\rceil.$$

Therefore

$$\min\{r \mid K_r \in F(G)\} \geq \left\lceil \frac{n}{\alpha(G)} \right\rceil. \quad \square$$

Theorem 5.7 Let G be a graph of order n with maximum degree Δ . Then

$$\max\{r \mid K_r \in F(G)\} \leq \left\lfloor \frac{1 + \sqrt{1 + 4n\Delta}}{2} \right\rfloor.$$

Proof: Let G be a graph of order n and with maximum degree Δ . Let $r = \max\{r \mid K_r \in F(G)\}$. By Theorem 2.5, $V(G)$ can be partitioned into

q independent sets S_1, S_2, \dots, S_q which are pairwise linked. Without loss of generality, assume that

$$k = |S_1| \leq |S_2| \leq \dots \leq |S_q|.$$

We claim that $k\Delta \geq q - 1$. Observe that the number of edges with one end vertex in S_1 is equal to

$$\sum_{x \in S_1} \deg(x).$$

Since S_1 is linked to all the other sets S_2, S_3, \dots, S_q , then

$$\sum_{x \in S_1} \deg(x) \geq q - 1.$$

Because $\forall x \in S_1, \deg(x) \leq \Delta$, then

$$\sum_{x \in S_1} \deg(x) \leq k\Delta.$$

It follows that

$$k\Delta \geq q - 1.$$

This proves our claim.

Then we have the following inequalities

$$n \geq qk \tag{1}$$

$$k\Delta \geq q - 1. \tag{2}$$

From these two inequalities, we get

$$n \geq q\left(\frac{q-1}{\Delta}\right)$$

by eliminating k . Thus

$$\begin{aligned} n\Delta &\geq q^2 - q \\ q^2 - q - n\Delta &\leq 0 \\ q &\leq \frac{1 + \sqrt{1 - 4(1)(-n\Delta)}}{2} \\ q &\leq \frac{1 + \sqrt{1 + 4n\Delta}}{2} \\ q &\leq \left\lfloor \frac{1 + \sqrt{4n\Delta + 1}}{2} \right\rfloor. \end{aligned}$$

Therefore

$$\max\{r \mid K_r \in F(G)\} \leq \left\lfloor \frac{1 + \sqrt{1 + 4n\Delta}}{2} \right\rfloor. \quad \square$$

Theorem 5.8 *Let G be a graph. Then*

$$\min\{r \mid K_r \in F(G)\} \geq \left\lceil \frac{\dim(G)}{2} \right\rceil.$$

Proof: By Theorem 5.2 and Lemma 4.5, we have

$$\min\{r \mid K_r \in F(G)\} \geq \chi(G)$$

and

$$\chi(G) \geq \frac{\dim(G)}{2}.$$

Hence

$$\min\{r \mid K_r \in F(G)\} \geq \frac{\dim(G)}{2}.$$

Therefore

$$\min\{r \mid K_r \in F(G)\} \geq \left\lceil \frac{\dim(G)}{2} \right\rceil. \quad \square$$

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