Bounding the Sum of Distances of the Vertices of a Tree

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The distance d(x, v) between the vertices x and y of a connected graph G is defined to be the length of a shortest path in G which joins x and y. It is clear that d(x,y) = d(y,x), $d(x,y) \ge 0$ and d(x,x) = 0.

and d(x,x) = 0. In this paper we will find a lower bound and an upper bound for the sum $\sum_{x,y \in V} d(x,y)$ of distances of vertices of a tree T = (V,E). To do this we shall investigate the family \Im of trees of order n and find out which trees in \Im will give extremum values for $\sum_{x,y \in V} d(x,y)$.

Lemma 1. $\sum_{i=1}^{n} i^2 = \frac{1}{6}n(n+1)(2n+1)$.

Proof We prove by induction on *n*. Let $S_n = \sum_{i=1}^n i^2$ and $f(n) = \frac{1}{6}n(n+1)(2n+1)$. Then $S_1 = 1 = f(1)$ and $S_2 = 5 = f(2)$. Suppose that n > 2

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 $\frac{1}{6}n(n+1)(2n+1)$. Then $S_1 = 1 - f(1)$ and $S_2 = 0 - f(1) - 0$ and $S_{n-1} = f(n-1)$. Now

$$S_n = S_{n-1} + n^2 = \frac{1}{6}(n-1)n(2n-1) + n^2$$
$$= \frac{1}{6}n(n+1)(2n+1).$$

This means that $S_n = f(n)$ for n > 2.

Lemma 2. Let T = (V,E) be a tree of order *n*. (1) If $T = K_{1,n-1}$ (star graph), then

$$\sum_{x,y\in V} d(x,y) = 2(n-1)^2.$$

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(2) If $T = P_n$ (path), then $\sum_{x,y \in V} d(x,y) = \frac{1}{3}n(n^2 - 1)$.

Proof. (1) Let $T = K_{1,n-1}$. Then consider the following cases:

<u>Case 1</u>. d(x,y) = 1. Then there are 2(n-1) terms of $\sum_{x,y \in V} d(x,y)$ for this case.

Case 2.
$$d(x,y) = 2$$
. Then there are $2\binom{n-1}{2} = (n-1)(n-2)$ terms of

 $\sum_{x,y\in V} d(x,y)$ for this case.

<u>Case 3</u>. d(x,y) > 2. Then there are no terms of $\sum_{x,y \in V} d(x,y)$ for this case Hence

$$\sum_{x,y \in V} d(x,y) = 2(n-1) + 2(n-1)(n-2) = 2(n-1)^2.$$

(2) For the second part we take $T = P_n = [1, 2, ..., n]$. Then consider the following cases:

- <u>Case 1</u>. d(x,y) = 1. Then there are 2(n-1) terms of $\sum_{x,y \in V} d(x,y)$ in this case.
- <u>Case 2</u>. d(x,y) = 2. Then [1, 2, 3], [2, 3, 4], ..., [n-2, n-1, n] are the subtrees of T whose lengths are each equal to 2. It follows that there are 2(n-2) terms of $\sum_{x,y \in V} d(x,y)$ in this case.
- <u>Case 3</u>. d(x,y) = 3. Then [1, 2, 3, 4], [2, 3, 4, 5], ..., [n-3, n-2, n-1, n]are the subtrees of T whose lengths are each equal to 3. Hence there are 2(n-3) terms of $\sum_{x,y\in V} d(x,y)$ in this case.

We continue in this fashion until we reach the last case where d(x,y) = n - 1. There are only 2 terms of $\sum_{x,y \in V} d(x,y)$ in this case. So for any tree T = (V, E) of order n,

$$\sum_{x,y\in V} d(x,y) = 2[(n-1) + 2(n-2) + 3(n-3) + \dots + (n-1)]$$

$$=2\sum_{i=1}^{n-1}i(n-i)=2\left[n\sum_{i=1}^{n-1}i-\sum_{i=1}^{n-1}i^{2}\right]$$

$$= 2 \left[\frac{1}{2} n^2 (n-1) - \sum_{i=1}^{n-1} i^2 \right].$$

Using Lemma 1, we can simplify the last expression above to obtain

$$\sum_{x,y \in V} d(x,y) = n^2 (n-1) - 2 \sum_{i=1}^{n-1} i^2$$
$$= n^2 (n-1) - \frac{1}{3} n(n-1)(2n-1)$$
$$= \frac{1}{3} n(n^2 - 1). \quad \Box$$

Theorem 3. Let \mathcal{T} be the family of trees T = (V,E) of order *n*. Then

$$\min\left\{\sum_{x,y\in V}d(x,y): T\in\mathfrak{I}\right\}=2(n-1)^2.$$

Proof. A tree of order n has n - 1 edges. Then in the sum $\sum_{x,y \in V} d(x,y)$ there are exactly 2(n - 1) terms each equal to 1 while the other nonzero terms are greater than or equal to 2. If $T = K_{1,n-1}$, then all the terms different from 0 and 1 are each equal to 2. If $T \neq K_{1,n-1}$, then there is at least one term which is equal to 3. Hence, the minimum value of $\sum_{x,y \in V} d(x,y)$ is attained only if $T = K_{1,n-1}$. Lemma 2 completes the proof of this theorem. \Box

Lemma 4. Let \mathcal{J} be the family of trees T = (V,E) of order *n*. The

sum

 $s(x) = \sum_{y \in V} d(x, y)$

is a maximum in \mathcal{T} if T is a path and x is an end-vertex.

Proof. Let T be of order *n*. For s(x) to be maximum, the n-1 edges of T must be constructed in such a way that d(x,y) is as large as possible for each $y \in V$. To do this, let $x_1, x_2, ..., x_n$ be the vertices of T where the indices correspond to the order in which T is contructed. Without loss of generality, let $x = x_1$. We begin with the edge $[x,x_2]$. For $d(x,x_3)$ to be maximum, we require that $[x_2,x_3] \in E$. Also, $[x_3,x_4]$ must be an edge of T for $d(x,x_4)$ to be maximum. Continuing in this manner, we

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can eventually require that $[x_{n-1}, x_n] \in E$. Hence it is necessary that P = $[x_1, x_2, x_3, ..., x_n]$ is a path in T. But P is of order n and it has n - 1 edges; hence T = P and x is an end-vertex. \Box

Theorem 5. Let \mathcal{J} be the family of trees T = (V, E) of order *n*. Then

$$\max\left\{\sum_{x,y\in V} d(x,y) : T \in \mathfrak{I}\right\} = \frac{1}{3}n(n^2-1).$$

Proof. First, we claim that $\sum_{x,y\in V} d(x,y)$ is a maximum in \Im if T is a path. We prove this claim by induction on n. For $1 \le n \le 3$, the property is immediate since every tree in this case is a path. Suppose now that the property is true for all trees with at most n - 1 vertices. Let z be an end-vertex of T. Then

$$\sum_{x,y \in \mathcal{V}} d(x,y) = 2s(z) + \sum_{x,y \in \mathcal{W}} d(x,y),$$

where $W = V - \{z\}$. Consider the right side of the last equation above. By Lemma 4, s(z) in the first term is maximum if T is a path and z is one of its end-vertices. The second term is also maximum by applying induction hypothesis to the subtree whose n - 1 vertices make up the set W. Thus $\sum_{x,y \in V} d(x,y)$ is a maximum. This completes the proof of the claim. By Lemma 2, the proof of the theorem is also completed. \Box

Corollary 6. For any tree T = (V,E) of order *n*,

$$2(n-1)^2 \leq \sum_{x,y \in V} d(x,y) \leq \frac{1}{3}n(n^2-1).$$

Proof. This corollary results from the combination of Theorem 3 and Theorem 5. \Box

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