Bounding the Sum of Distances of the Vertices of a Tree

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The distance $d(x,y)$ between the vertices x and y of a connected f_{ne} is defined to be the length of a shortest path f_{in} edge of f_{in} which joins *x* and *y*. It is clear that $a(x,y) = a(y,x)$, $d(x,y) \ge 0$

and $d(x,x) = 0$.
In this paper we will find a lower bound and an upper bound for sum $\sum_{x,y\in V} d(x,y)$ of distances of vertices of a tree $T = (V,E)$. To do this we shall investigate the family 3 of trees of order *n* and find out which trees in 3 will give extremum values for $\sum_{x,y\in V} d(x,y)$.

Lemma 1. $\sum_{i=1}^{n} i^2 = \frac{1}{6}n(n+1)(2n+1)$.

Proof We prove by induction on *n*. Let $S_n = \sum_{i=1}^{n} i^2$ and $f(n) =$ $\frac{1}{6}n(n+1)(2n+1)$. Then $S_1 = 1 = f(1)$ and $S_2 = 5 = f(2)$. Suppose that $n > 2$

and $S_{n-1} = f(n-1)$. Now

$$
S_n = S_{n-1} + n^2 = \frac{1}{6}(n-1)n(2n-1) + n^2
$$

= $\frac{1}{6}n(n+1)(2n+1)$.

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This means that $S_n = f(n)$ for $n > 2$. \Box

 $graph$, then \Box **Lemma 2**. Let $T = (V,E)$ be a tree of order *n*. (1) If $T = K_{1,n-1}$ (star

$$
\sum_{x,y\in V} d(x,y) = 2(n-1)^2.
$$

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(2) If $T = P_n$ (path), then $\sum_{x,y \in V} d(x,y) = \frac{1}{3}n(n^2 - 1)$.

Proof. (1) Let $T = K_{1,n-1}$. Then consider the following cases:

<u>Case 1</u>. $d(x,y) = 1$. Then there are $2(n-1)$ terms of $\sum_{x,y\in V} d(x,y)$ for this case.

Case 2.
$$
d(x,y) = 2
$$
. Then there are $2{n-1 \choose 2} = (n-1)(n-2)$ terms of

 $\sum_{x,y\in V} d(x,y)$ for this case.

Case 3. $d(x,y) > 2$. Then there are no terms of $\sum_{x,y\in V} d(x,y)$ for this case Hence

$$
\sum_{x,y \in V} d(x,y) = 2(n-1) + 2(n-1)(n-2) = 2(n-1)^2.
$$

(2) For the second part we take $T = P_n = [1, 2, ..., n]$. Then consider the following cases:

- <u>Case 1</u>. $d(x,y) = 1$. Then there are $2(n-1)$ terms of $\sum_{x,y\in V} d(x,y)$ in this case.
- Case 2. $d(x,y) = 2$. Then [1, 2, 3], [2, 3, 4], ..., $[n-2, n-1, n]$ are the subtrees of T whose lengths are each equal to 2. It follows that there are $2(n-2)$ terms of $\sum_{x,y\in V} d(x,y)$ in this case.
- Case 3. $d(x,y) = 3$. Then [1, 2, 3, 4], [2, 3, 4, 5], ..., $[n-3, n-2, n-1, n]$ are the subtrees of T whose lengths are each equal to 3. Hence there are $2(n-3)$ terms of $\sum_{x,y\in V} d(x,y)$ in this case.

We continue in this fashion until we reach the last case where $d(x, y)$ $= n - 1$. There are only 2 terms of $\sum_{x,y \in V} d(x,y)$ in this case. So for any tree $T = (V, E)$ of order *n*,

$$
\sum_{x,y \in V} d(x,y) = 2[(n-1)+2(n-2)+3(n-3)+...+(n-1)]
$$

$$
=2\sum_{i=1}^{n-1}i(n-i)=2\left[n\sum_{i=1}^{n-1}i-\sum_{i=1}^{n-1}i^{2}\right]
$$

$$
=2\bigg[\tfrac{1}{2}n^2(n-1)-\sum_{i=1}^{n-1}i^2\bigg].
$$

Using Lemma 1, we can simplify the last expression above to obtain

$$
\sum_{x,y \in V} d(x,y) = n^2(n-1) - 2\sum_{i=1}^{n-1} i^2
$$

= $n^2(n-1) - \frac{1}{3}n(n-1)(2n-1)$
= $\frac{1}{3}n(n^2-1)$. \square

'Then **Theorem 3.** Let \mathcal{T} be the family of trees $T = (V, E)$ of order

$$
\min\left\{\sum_{x,y\in V}d(x,y): \mathrm{T}\in \mathfrak{T}\right\}=2(n-1)^2.
$$

Proof. A tree of order *n* has $n - 1$ edges. Then in the sum $\sum_{x,y\in V} d(x,y)$ there are exactly $2(n - 1)$ terms each equal to 1 while the other nonzero terms are greater than or equal to 2. If $T = K_{1,n-1}$, then all the terms different from 0 and 1 are each equal to 2. If $T \neq K_{1,n-1}$, then there is at least one term which is equal to 3. Hence, the minimum value of $\sum_{x,y\in V} d(x,y)$ is attained only if T = K_{1,n-1}. Lemma 2 completes the proof of this theorem. \Box

Lemma 4. Let \mathcal{T} be the family of trees $T = (V,E)$ of order *n*. The

sum

$$
s(x) = \sum_{y \in V} d(x, y)
$$

is a maximum in $\mathcal I$ if T is a path and x is an end-vertex.

Proof. Let T be of order *n*. For $s(x)$ to be maximum, the $n-1$ edges of T must be constructed in such a way that $d(x,y)$ is as large as possible for each $y \in V$. To do this, let $x_1, x_2, ..., x_n$ be the vertices of T where the indices correspond to the order in which T is contructed. Without loss of generality, let $x = x_1$. We begin with the edge [x,x₂]. For $d(x, x_3)$ to be maximum, we require that $[x_2, x_3] \in E$. Also, $[x_3, x_4]$ must be an edge of T for $d(x,x_4)$ to be maximum. Continuing in this manner, we

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 $\begin{bmatrix} \n \text{can eventually require that} & [x_{n-1}, x_n] \in E. \n \end{bmatrix}$ Hence it is necessary that **P** $=$ $[x_1, x_2, x_3, \ldots, x_n]$ is a path in T. But P is of order *n* and it has $n-1$ edges; hence $T = P$ and *x* is an end-vertex.

Then **Theorem 5.** Let \mathcal{T} be the family of trees $T = (V, E)$ of order *n*.

$$
\max \left\{ \sum_{x,y \in V} d(x,y) : T \in \mathfrak{I} \right\} = \frac{1}{3} n(n^2-1).
$$

Proof. First, we claim that $\sum_{x,y\in V} d(x,y)$ is a maximum in \Im if T is a path. We prove this claim by induction on *n*. For $1 \le n \le 3$, the property is immediate since every tree in this case is a path. Suppose now that the property is true for all trees with at most $n-1$ vertices. Let *z* be an endvertex of T. Then

$$
\sum_{x,y \in V} d(x,y) = 2s(z) + \sum_{x,y \in W} d(x,y),
$$

where $W = V - \{z\}$. Consider the right side of the last equation above. By Lemma 4, $s(z)$ in the first term is maximum if T is a path and z is one of its end-vertices. The second term is also maximum by applying induction hypothesis to the subtree whose $n-1$ vertices make up the set W. Thus $\sum_{x,y\in V} d(x,y)$ is a maximum. This completes the proof of the claim. By Lemma 2, the proof of the theorem is also completed. \mathcal{C} .

Corollary 6. For any tree $T = (V,E)$ of order *n*,

$$
2(n-1)^{2} \leq \sum_{x,y \in V} d(x,y) \leq \frac{1}{3} n(n^{2}-1).
$$

Proof. This corollary results from the combination of Theorem 3 and Theorem 5. \Box

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